Topic Review for Finals

# Gale-Shapley and Stable Matching

Full list of preferences, no ties. Gale-Shapley returns a perfect and stable matching. Men are matched to their best valid partners. Women are matched to their worst valid partners. The runtime is , being the number of men or the number of women, which are equal. Important things to know for proving stuff is that men go down their list of preferences, and women’s offers get better.

# Greedy Algorithms

## Unweighted Interval Scheduling

Input is a set of requests (time intervals). Goal is to return a set of non-overlapping intervals of maximum cardinality possible. The size of each interval doesn’t matter, we just want as many non-overlapping intervals as possible. Selecting earliest start time then deleting conflicts doesn’t work, selecting by increasing number of overlaps then deleting conflicts doesn’t work, selecting by increasing interval length and deleting conflicts doesn’t work. The algorithm is to select by **earliest finish time**! You can prove this by the exchange argument and induction.

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## Minimum Spanning Tree

### Kruskal

Sort all edges in increasing order of cost. Pick the smallest edge. If it forms a cycle with the tree we have so far, throw it out. Else, include it. Repeat that until all nodes are included.

Kruskal can be in runtime complexity if using the disjoint set data structure.

### Prim

Start from random node. Check outgoing edges of the tree we have so far, pick the one with the lowest cost. If it creates a cycle, throw it out. Else, keep it. Repeat that until all nodes are included. If all weights in a graph are unique, then there is a unique minimum spanning tree, so all MST algorithms returns the same thing. Just a useful thing to know, the time complexity of DFS is .

Prim can be in runtime complexity if using priority heaps to store edges.

### Borůvka

Initialize all vertices to be a component. While there is more than one component (i.e. no spanning tree yet), find the cheapest edge for each component and add it to our current trees (this is done in parallel for all components). Requires distinct edge weights. Because imagine a triangle with the same weight on each edge. Each one can choose the edge to the left and we would end up with a cycle.

Borůvka can be in runtime complexity.

### Cut Property

It is a way to prove MST algorithms. The claim could be that given node and edge that is the cheapest between the set it’s in and , is included in every MST. The cut property says that assume the optimal solution doesn’t have , but has that connects to . There must be a path from in to in . Then adding to this optimal MST creates a cycle. Just remove and add in to create a new solution. This solution is cheaper or not more expensive than .

## Huffman Coding

Not done in homework or practice sets. Do not need to know proof. Just need to know reduction to other problems.

We want binary representations for each letter in a lexicon, want shorter strings for more frequently used letters and longer strings for the less frequently used ones. Also! (This is important.) We don’t want a letter’s coding to be the prefix to another letter’s coding, i.e. if a’s code is 10, then other letter’s code can’t start with 01, like c can’t be 101. Another objective is to have the average length of all codes as short as possible.

The answer is to use trees, with the nodes as the letters traversed. This way, one letter’s code can’t be another letter’s prefix. A more frequent letter appears shallower (upper) in the tree. What they want us to know if we’re *given* a tree, just know to assign the letters to increasingly deeper leaves in decreasing frequency. They also want us to know how to construct a tree given a table of letters and their frequencies. This is done by picking the two least frequent letters or components and making them siblings in the tree. Then take the sum of their frequencies and make the frequency on their parent that. Repeat.

Huffman coding’s runtime complexity is .

# Dynamic Programming

## Weighted Interval Scheduling

Interval scheduling, but with value associated to each interval. Want compatible set of intervals of greatest value. Sort by finishing time first! This is important! All recurrence schemes are useless without some metric of order. Interval could be or not be included. The same problem has to be solved on the previous intervals. Define to be the latest interval that doesn’t conflict with interval . If it is included, then the latest interval that’s also there is the latest that doesn’t overlap. If it isn’t included, then the latest one would be the previous in the ordered sequence. The recurrence scheme is , base case is .

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## Segmented Least Squares

Not done in homework or practice sets. Do not need to know proof. Just need to know reduction to other problems.

Given a bunch of points and a real number , you want to fit them into a number of lines of best fits. For each line of best fit, there’s an error of fitting the points to the line. The goal is minimizing the penalty , being the segments. The way to think about this is that if the last segment of the optimal solution is , then the value of the optimal solution is . The recurrence is thus . The base case is .

Segmented least squares’ runtime complexity is .

## Sequence Alignment

Not done in homework or practice sets. Do not need to know proof. Just need to know reduction to other problems.

Given two finite-length strings and from the same alphabet. Matchings, if represented as lines connecting matching letters between the two strings, cannot intersect. I.e. A matching is a pair matching with such that if there’s two matchings, and , that if then . The main objective is that given meaning the cost of changing an to a and meaning the cost of inserting or deleting an , we want to minimize costs of editing for and to match. Base cases are and . Recurrence scheme: because there’s either a matching between and , or is not matched with anything in , or is not matching with anything in .

Sequence alignment’s runtime complexity is .

## Bellman-Ford

Given a graph and two nodes from it, and , find the (value of) the shortest path from to . Having negative cycles in graphs is impossible to get a shortest path, because you will just loop on it forever instead of moving forward to the goal. Dijkstra’s algorithm will not work on graphs with negative weight edges (not even negative cycles, just weights), because it relies on the simple fact that adding a path can never make a path shorter; when there’s negative edges, a myopic cheap edge might lure you down a path of positive weights while an expensive one could hide negative edges behind. The closed node will never have to be reopened in an all-positive graph. The base cases for Bellman-Ford are and for all nodes other than . The recurrence scheme is , where the two arguments here mean the it’s the cost of the shortest path from to using at most edges. The recurrence scheme is as such because you either use at most edges, or you use edges and the node immediately before is .

Bellman-Ford’s runtime complexity is .

# Divide and Conquer

## The Master Method

is the size of the problem.

is the number of subproblems in the recursion.

is the size of each subproblem.

is the cost of the work done outside the recursive calls, which includes the cost of dividing the problem and the cost of merging the solution to the subproblems.

### Case

### Case

### Case 3

## Integer Multiplication

Multiplying two numbers and , each of places. The solution is as follows. Write as , write similarly as . Compute . Compute , . The answer is .

The runtime complexity of Karatsuba multiplication is because is , is , and is .

## Convolution

Only to recognize a convolution problem, not details of how to run the algorithm.

Given two vectors and , want

# Randomized Algorithms

## Linear-Time Median Finding

We pick a pivot at random at each recursive step. We want the pivot to be “good,” that at least elements in the array is less than it and at least elements is greater than it. The expected number of of repeats to get such a pivot in a recursive step is . Using the Master method on this produces , as each subproblem is at the biggest .

This is runtime complexity.

## Hashing

The goal of hashing is to have a good function such that it stores at position and have very few collisions. Formally, a universal hash function is one where the probability of a collisions is . Making a universal hash function involves picking a prime number .

Probability of a collision here is .

## Prime Testing

The probability of a number of digits being prime is . Fermat’s theorem is that if is prime, then , . So we use the contrapositive, that if , then is not prime.

# Network Flow

## Ford-Fulkerson

Given a graph with directed edges, each with a capacity , we want to send as much flow as possible from to . Forward flow subtracts from the capacity. “Backward flow” adds to the capacity. At each step in the algorithm, choose any path from to and update the residual graph accordingly. It terminates when there are no more such paths. It is important that all capacities and flows are integers! We need to say that in proofs, too.

Ford-Fulkerson’s runtime complexity is , where is the sum of the capacities leaving and is the number of edges.

## Min Cut

The max-flow min-cut theorem says the maximum flow is equal to the total capacities of the edges in the minimum cut, i.e. the smallest total weight of the edges which if removed would disconnect the source from the sink.

The capacity of a cut and such that and is the sum of the capacities of edges leaving . There is no subtracting capacities of edges entering A! To get the minimum cut, run Ford-Fulkerson and get the final residual graph, then is the set of all nodes that you can reach from .

If a flow’s value equals a cut’s capacity, then that is the max flow and that is the min cut. Because that’s the biggest flow for that cut, and that is the smallest cut for that flow.

## Applications

### Disjoint Paths

Given a directed or undirected graph, we want to find the maximum number of edge-disjoint paths in it. There are edge-disjoint graphs in a directed graph if and only if the max flow from to (minimum number of edges whose removal separates from ) is .

### Project Selection

We’re given a set of projects, each with a profit, and each of which might have prerequisites, and each of those has a cost (negative profit). What is the set of projects to do (along with their required prerequisites) such that profit is maximized? The solution is to find a min cut in a graph constructed as such: connect to all projects that generate a profit (don’t consider the prerequisites, just itself) and give them their profit as edge capacity; connect projects to their prerequisites with edges of infinite capacity (this is to make sure a cut doesn’t cross these edges, so as to preserve prerequisite structure); and connect all negative profit projects to with capacity (so, a positive number).

# NP-Completeness

## SAT problem

A bunch of conjunctions (ands) of clauses, each clause is a bunch of disjunctions (ors). Is there an assignment for all the variables to satisfy the whole thing?

## Independent Set

Given a graph, is there a set of nodes of size that aren’t connected to each other?

## Vertex Cover

Given a graph, is there a set of nodes whose incident edges touch all nodes in the graph?

## Set Cover

Given a set of sets, is there a subset of the larger set of size such that their union is the union of all the sets in the larger set?

## Hamiltonian Path

A Hamiltonian path is a path between two vertices of a graph that touches every vertex. If the two vertices are adjacent, then it’s a Hamiltonian cycle.

## Traveling Salesman Problem

Goal is to get a Hamiltonian cycle, but of the cheapest edge cost.

# Computability

## Halting Problem

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## Co-Halting Problem

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## Accept Problem

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# Approximation Algorithms

## Greedy Methods

### Knapsack

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## Linear Programming

### Vertex Cover

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